

### On a Theorem in Graph Theory

An error in the proof of a theorem given in Ref. 1 is pointed out and a correct proof is provided.

Seshu and Reed<sup>1</sup> have proved the following theorem in their book<sup>1</sup>:

If  $G$  is a connected graph of  $v$  vertices and  $G_s$  is a subgraph of  $G$  with  $v-1$  elements and containing no circuits, then  $G_s$  is a tree of  $G$ .

In this correspondence, we first point out an error in the proof of the above theorem and then provide a correct proof.

This theorem will be true if it can be shown that:

- (1)  $G_s$  contains  $v$  vertices, and
- (2)  $G_s$  is connected.

In their proof, Seshu and Reed first state that  $G_s$  contains all the  $v$  vertices of  $G$  and then establish (2). They then argue that since  $G_s$  is connected, contains  $v-1$  elements and no circuits, it is its own tree and hence contains  $v$  vertices. Thus they establish (2) assuming (1) to be true and then establish (1) assuming (2) to be true. Since (1) and (2) should be proved independently, the proof as given in Ref. 1 is not valid. The following comments are also in order:

(i) by definition [def. 1-4 in Ref. 1], a subgraph  $G_s$  of a connected graph  $G$  consists of a subset of edges (and hence also the vertices at which these edges are incident) of  $G$ . Thus according to this definition  $G_s$  need not, in general, contain all the vertices of  $G$ . Hence equation (2-5) used in Ref. 1 is not correct.

(ii) Suppose we define that a subgraph  $G_s$  of a graph  $G$  consists of all the vertices of  $G$  and a subset of edges of  $G$ . Then equation (2-5) is correct. However, in such a case, the statements 'Now  $G_s$  is its own tree and contains  $v-1$  elements. Hence  $G_s$  contains  $v$  vertices . . . are not necessary. This latter definition of a subgraph is not the usual one. Such a definition, though convenient for the proof of this theorem, will lead to difficulties in the development of other results of graph theory.

Thus a correct proof of the theorem is called for. This is detailed as follows:

Let  $G_s$  consist of  $p$  maximal connected subgraphs. Let  $s_1, s_2, s_3, \dots, s_p$  be the number of vertices in  $s_i$ . Since  $s_i$  is connected and contains no circuits  $s_i$  is its own tree. Hence  $s_i$  contains  $v_i - 1$  elements. Thus  $G_s$  contains  $v^* - p$  elements where

$$v^* = \sum_{i=1}^p v_i$$

Since, by hypothesis,  $G_s$  contains  $v-1$  elements, we get

$$v^* - p = v - 1 \quad (1)$$

Hence

$$v^* - v = p - 1 \quad (2)$$

However,  $v^* \leq v$  and hence we get from (2)

$$p^* \leq 1 \quad (3)$$

But

$$p \geq 1 \quad (4)$$

Hence from (3) and (4) it follows that  $p = 1$ . Thus  $G_p$  is connected.

Since  $p = 1$  we get from (1)

$$v^* = v$$

Hence  $G_p$  contains all the  $v$  vertices. Thus follows the proof of the theorem.

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#### Reference

1. *S. Seshu and M. B. Reed* Linear Graphs and Electrical Networks, p. 26, Addison Wesley Publishing Company, Reading, Massachusetts, U.S.A.

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