

Inference Rules: Propositional Calculus

$$\frac{a \quad b}{a \wedge b} \{\wedge I\}$$

$$\frac{a \wedge b}{a} \{\wedge E_L\}$$

$$\frac{a \wedge b}{b} \{\wedge E_R\}$$

$$\frac{a}{a \vee b} \{\vee I_L\}$$

$$\frac{b}{a \vee b} \{\vee I_R\}$$

$$\frac{a \vee b \quad a \vdash c \quad b \vdash c}{c} \{\vee E\}$$

$$\frac{a \vdash b}{a \rightarrow b} \{\rightarrow I\}$$

$$\frac{a \quad a \rightarrow b}{b} \{\rightarrow E\}$$

$$\frac{a}{a} \{ID\}$$

$$\frac{\text{False}}{a} \{CTR\}$$

$$\frac{\neg a \vdash \text{False}}{a} \{RAA\}$$

Some Theorems in Rule Form

$$\frac{a \wedge b}{b \wedge a} \{\wedge\text{Comm}\}$$

And Commutes

$$\frac{a \vee b}{b \vee a} \{\vee\text{Comm}\}$$

Or Commutes

$$\frac{}{a \vee (\neg a)} \{\text{noMiddle}\}$$

Law of Excluded Middle

$$\frac{a \rightarrow b \quad b \rightarrow c}{a \rightarrow c} \{\rightarrow\text{Chain}\}$$

Implication Chain Rule

$$\frac{\neg(a \vee b) \quad \neg(b \vee a)}{\neg(\neg(\vee)\text{Comm})}$$

Not Or Commutes

$$\frac{a \rightarrow b \quad \neg b}{\neg a} \{\text{modTo}\}$$

Modus Tollens

$$\frac{a \rightarrow b}{(\neg b) \rightarrow (\neg a)} \{\text{conPos}_F\}$$

Contrapositive Fwd

$$\frac{}{\text{False}} \{\neg\&\}$$

NeverBoth

$$\frac{a \rightarrow b}{(\neg a) \vee b} \{\rightarrow_F\}$$

Implication Fwd

$$\frac{(\neg a) \vee b}{a \rightarrow b} \{\rightarrow_B\}$$

Implication Bkw

More Theorems in Rule Form

$$\frac{\neg(a \vee b)}{(\neg a) \wedge (\neg b)} \quad \{DeM_{\vee F}\}$$

DeMorgan Or Fwd

$$\frac{(\neg a) \wedge (\neg b)}{\neg(a \vee b)} \quad \{DeM_{\vee B}\}$$

DeMorgan Or Bkw

$$\frac{\neg(a \wedge b)}{(\neg a) \vee (\neg b)} \quad \{DeM_{\wedge F}\}$$

DeMorgan And Fwd

$$\frac{(\neg a) \vee (\neg b)}{\neg(a \wedge b)} \quad \{DeM_{\wedge B}\}$$

DeMorgan And Bkw

$$\frac{a \vee b \quad \neg a}{b} \quad \{disjSyl/\}$$

Disjunctive Syllogism

$$\frac{\neg(\neg a)}{a} \quad \{\neg \neg_F\}$$

Double Negation Fwd

$$\frac{a}{\neg(\neg a)} \quad \{\neg \neg_B\}$$

Double Negation Bkw

| | |
|--|--|
| $a \wedge \text{False} = \text{False}$ | { $\wedge \text{null}$ } |
| $a \vee \text{True} = \text{True}$ | { $\vee \text{null}$ } |
| $a \wedge \text{True} = a$ | { $\wedge \text{identity}$ } |
| $a \vee \text{False} = a$ | { $\vee \text{identity}$ } |
| $a \wedge a = a$ | { $\wedge \text{idempotent}$ } |
| $a \vee a = a$ | { $\vee \text{idempotent}$ } |
| $a \wedge b = b \wedge a$ | { $\wedge \text{commutative}$ } |
| $a \vee b = b \vee a$ | { $\vee \text{commutative}$ } |
| $(a \wedge b) \wedge c = a \wedge (b \wedge c)$ | { $\wedge \text{associative}$ } |
| $(a \vee b) \vee c = a \vee (b \vee c)$ | { $\vee \text{associative}$ } |
| $a \wedge (b \vee c) = (a \wedge b) \vee (a \wedge c)$ | { $\wedge \text{distributes over } \vee$ } |
| $a \vee (b \wedge c) = (a \vee b) \wedge (a \vee c)$ | { $\vee \text{distributes over } \wedge$ } |
| $\neg(a \wedge b) = (\neg a) \vee (\neg b)$ | {DeMorgan's law \wedge } |
| $\neg(a \vee b) = (\neg a) \wedge (\neg b)$ | {DeMorgan's law \vee } |
| $\neg \text{True} = \text{False}$ | {negate True} |
| $\neg \text{False} = \text{True}$ | {negate False} |
| $(a \wedge (\neg a)) = \text{False}$ | { $\wedge \text{complement}$ } |
| $(a \vee (\neg a)) = \text{True}$ | { $\vee \text{complement}$ } |
| $\neg(\neg a) = a$ | {double negation} |
| $(a \wedge b) \rightarrow c = a \rightarrow (b \rightarrow c)$ | {Currying} |
| $a \rightarrow b = (\neg a) \vee b$ | {implication} |
| $a \rightarrow b = (\neg b) \rightarrow (\neg a)$ | {contrapositive} |

Axioms

Some Equations of Boolean Algebra

Theorems

$$(a \wedge b) \vee b = b \quad \{\vee \text{absorption}\}$$

$$(a \vee b) \wedge b = b \quad \{\wedge \text{absorption}\}$$

$$(a \vee b) \rightarrow c = (a \rightarrow c) \wedge (b \rightarrow c) \quad \{\vee \text{imp}\}$$